Dynamics and Equilibrium

Sergiu Hart

Presidential Address, GAMES 2008 (July 2008)

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DYNAMICS AND EQUILIBRIUM

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Papers

Papers

- Hart and Mas-Colell, Econometrica 2000
- Hart and Mas-Colell, J Econ Theory 2001
- Hart and Mas-Colell, Amer Econ Rev 2003
- Hart, Econometrica 2005
- Hart and Mas-Colell, Games Econ Behav 2006
- Hart and Mansour, Games Econ Behav 2009?
- Hart, Center for Rationality DP 2008

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http://www.ma.huji.ac.il/hart

John Nash, Ph.D. Dissertation, Princeton 1950

EQUILIBRIUM POINT:

John Nash, Ph.D. Dissertation, Princeton 1950

EQUILIBRIUM POINT:

"Each player's strategy is optimal against those of the others."

John Nash, Ph.D. Dissertation, Princeton 1950

FACT

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There are no general, natural dynamics leading to Nash equilibrium

"general"

FACT

There are no general, natural dynamics leading to Nash equilibrium

"general": in all games

FACT

There are no general, natural dynamics leading to Nash equilibrium

"general": in all games rather than: in specific classes of games

FACT

- "general": in all games rather than: in specific classes of games:
 - two-person zero-sum games
 - two-person potential games
 - supermodular games
 - **.** . . .

FACT

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"leading to Nash equilibrium"

FACT

There are no general, natural dynamics leading to Nash equilibrium

"leading to Nash equilibrium": at a Nash equilibrium (or close to it) from some time on

FACT

FACT

There are no general, natural dynamics leading to Nash equilibrium

"natural"

FACT

There are no general, natural dynamics leading to Nash equilibrium

"natural":

FACT

- "natural":
 - adaptive (reacting, improving, ...)

FACT

- "natural":
 - adaptive (reacting, improving, ...)
 - simple and efficient

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 - computation (performed at each step)

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- "natural":
 - adaptive (reacting, improving, ...)
 - simple and efficient:
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bounded rationality

Dynamics that are **NOT** "natural":

 exhaustive search (deterministic or stochastic)

- exhaustive search (deterministic or stochastic)
- using a mediator

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- broadcasting the private information and then performing joint computation

- exhaustive search (deterministic or stochastic)
- using a mediator
- broadcasting the private information and then performing joint computation
- fully rational learning
 (prior beliefs on the strategies of the opponents, Bayesian updating, optimization)

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Hart and Mas-Colell, AER 2003

UNCOUPLED DYNAMICS:

Each player knows only his own payoff (utility) function

(does *not* know the payoff functions of the other players)

(privacy-preserving, decentralized, distributed ...)

Hart and Mas-Colell, AER 2003

Games

N-person game in strategic (normal) form:

Players

$$i=1,2,...,N$$

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For each player i: Actions

$$oldsymbol{a^i}$$
 in $oldsymbol{A^i}$

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N-person game in strategic (normal) form:

Players

$$i=1,2,...,N$$

For each player i: Actions

$$a^i$$
 in A^i

For each player i: Payoffs (utilities)

$$\mathbf{u}^{i}(a) \equiv \mathbf{u}^{i}(a^{1}, a^{2}, ..., a^{N})$$

Time

$$t = 1, 2, ...$$

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ullet At period t each player i chooses an action a_t^i in A^i

Time

$$t = 1, 2, ...$$

At period t each player i chooses an action

$$oldsymbol{a_t^i}$$
 in A^i

according to a probability distribution

$$oldsymbol{\sigma_t^i}$$
 in $\Delta(A^i)$

Fix the set of players 1, 2, ..., N and their action spaces $A^1, A^2, ..., A^N$

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A general dynamic:

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A general dynamic:

$$\sigma_t^i \equiv \sigma_t^i$$
 (HISTORY; GAME)

Fix the set of players 1, 2, ..., N and their action spaces $A^1, A^2, ..., A^N$

A general dynamic:

$$egin{aligned} \sigma_t^i &\equiv \sigma_t^i \ (ext{ HISTORY} \; ; \; ext{GAME} \) \ &\equiv \sigma_t^i \ (ext{ HISTORY} \; ; \; u^1,...,u^i,...,u^N \) \end{aligned}$$

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$$\sigma_t^i \equiv \sigma_t^i \, (\, {\sf HISTORY} \, ; \, {\color{red} u^i} \,)$$

Simplest uncoupled dynamics

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where $a_{t-1}=(a_{t-1}^1,a_{t-1}^2,...,a_{t-1}^N)\in A$ are the actions of all the players in the previous period

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Only last period matters ("1-recall")

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- Only last period matters ("1-recall")
- Time t does not matter ("stationary")

Impossibility

Impossibility

Theorem. There are **NO** uncoupled dynamics with 1-recall

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that yield almost sure convergence of play to pure Nash equilibria of the stage game in all games where such equilibria exist.

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Consider the following two-person game, which has a unique pure Nash equilibrium

	C1	C2	C3
R1	1,0	0,1	1,0
R2	0,1	1,0	1,0
R3	0,1	0,1	1,1

Consider the following two-person game, which has a unique pure Nash equilibrium (R3,C3)

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Assume *by way of contradiction* that we are given an uncoupled, 1-recall, stationary dynamic that yields almost sure convergence to pure Nash equilibria when these exist

• Suppose the play at time t-1 is (R1,C1)

	C1	C2	C3
R1	1,0	0,1	1,0
R2	0,1	1,0	1,0
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- Suppose the play at time t-1 is (R1,C1)
- ROWENA is best replying at (R1,C1)

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- Suppose the play at time t-1 is (R1,C1)
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- \Rightarrow ROWENA will play R1 also at t Proof:
 - Change the payoff function of COLIN so that (R1,C1) is the unique pure Nash eq.

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 - By uncoupledness, the same holds in the original game

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- ightharpoonup Rowena will play the same action at t

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R1	1,0 ↔	0,1	1,0
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Similarly for COLIN:

	C1	C2	C3
R1	1,0 ↔	0,1 🙏	1,0 ↔
R2	0,1 🙏	1,0 ↔	1,0 ↔
R3	0,1	0,1 🙏	1,1

Similarly for COLIN:

A player who is best replying cannot switch

	C1	C2	C3
R1	1,0 ↔	0,1 🙏	1,0 ↔
R2	0,1 🙏	1,0 ↔	1,0 ↔
R3	0,1 🙏	0,1 🚺	1,1

⇒ (R3,C3) cannot be reached

Similarly for COLIN:

	C1	C2	C3
R1	1,0 ↔	0,1 🙏	1,0 ↔
R2	0,1 🙏	1,0 ↔	1,0 ↔
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Theorem. THERE EXIST uncoupled dynamics with 2-RECALL

$$\sigma_t^i \equiv f^i(a_{t-2},a_{t-1};u^i)$$

that yield almost sure convergence of play to pure Nash equilibria of the stage game in every game where such equilibria exist.

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IF.

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- ullet Player i best replied: $a^i \in \mathrm{BR}^i(a^{-i};u^i)$

THEN: At t player i plays a^i again: $a^i_t = a^i$

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THEN: At t player i plays a^i again: $a^i_t = a^i$

ELSE: At t player i randomizes uniformly over A^i

"Good":

"Good":

simple

"Good":

simple

"Bad":

"Good":

simple

"Bad":

exhaustive search

"Good":

simple

"Bad":

- exhaustive search
- all players must use it

"Good":

simple

"Bad":

- exhaustive search
- all players must use it
- takes a long time

FACT

- "natural":
 - adaptive
 - simple and efficient:
 - computation
 - time
 - information

FACT

- "natural":
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 - information: uncoupledness √

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 - **computation**: finite recall √
 - time
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FACT

- "natural":
 - adaptive
 - simple and efficient:
 - **computation**: finite recall √
 - time to reach equilibrium ?
 - information: uncoupledness √

HOW LONG TO EQUILIBRIUM?

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Estimate the number of time periods it takes until a Nash equilibrium is reached

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- How?
- An uncoupled dynamic



A distributed computational procedure

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Estimate the number of time periods it takes until a Nash equilibrium is reached

- How?
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 \approx

A distributed computational procedure

■ ⇒ COMMUNICATION COMPLEXITY

Distributed computational procedures

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 - START: Each participant has some private information

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Yao 1979, Kushilevitz and Nisan 1997

How Long to Equilibrium

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Uncoupled dynamics leading to Nash equilibria

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Conitzer and Sandholm 2004

An uncoupled dynamic leading to Nash equilibria is TIME-EFFICIENT if

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Theorem. There are **NO TIME-EFFICIENT** uncoupled dynamics that reach a pure Nash equilibrium in all games where such equilibria exist.

• An uncoupled dynamic leading to Nash equilibria is TIME-EFFICIENT if its COMMUNICATION COMPLEXITY is POLYNOMIAL in the number of players (rather than: exponential)

Theorem. There are **NO TIME-EFFICIENT** uncoupled dynamics that reach a pure Nash equilibrium in all games where such equilibria exist.

Hart and Mansour, GEB 2009 (?)

Intuition:

- Intuition:
 - different games have different equilibria

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- the dynamic procedure must distinguish between them

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- different games have different equilibria
- the dynamic procedure must distinguish between them
- no single player can do so by himself

FACT

There are No general, natural dynamics leading to Nash equilibrium

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RESULT

There CANNOT BE general, natural dynamics leading to Nash equilibrium

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Perhaps we are asking too much?

RESULT

There CANNOT BE general, natural dynamics leading to Nash equilibrium

- Perhaps we are asking too much?
- For instance, the size of the data (the payoff functions) is exponential rather than polynomial in the number of players

CORRELATED EQUILIBRIUM

Aumann, JME 1974

CORRELATED EQUILIBRIUM:

Nash equilibrium when players receive payoff-irrelevant information before playing the game

Aumann, JME 1974

A Correlated Equilibrium is a Nash equilibrium when players receive payoff-irrelevant signals before playing the game

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- Independent signals Nash equilibrium
- Public signals ("sunspots")

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 convex combinations of Nash equilibria

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 convex combinations of Nash equilibria
- Butterflies play the Chicken Game ("Speckled Wood" Pararge aegeria)

"Chicken" game

LEAVE STAY

LEAVE

STAY

5,5	3,6
6,3	0,0

"Chicken" game

	LEAVE	STAY
LEAVE	5,5	3,6
STAY	6,3	0,0

a Nash equilibrium

"Chicken" game

LEAVE STAY

LEAVE

STAY

5, 5	3,6

 $6,3 \quad | \quad 0,0$

another Nash equilibrium

"Chicken" game

LEAVE STAY

LEAVE 5,5

STAY

5, 5	3,6
6,3	0,0

0	1/2
1/2	0

a (publicly) correlated equilibrium

"Chicken" game

	LEAVE	STAY	
LEAVE	5,5	3,6	
STAY	6,3	0,0	

L
$$1/3$$
 $1/3$ S $1/3$ 0

another correlated equilibrium

- after signal L play LEAVE
- after signal s play STAY

A Correlated Equilibrium is a Nash equilibrium when the players receive payoff-irrelevant signals before playing the game (Aumann 1974)

- Independent signals Nash equilibrium
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 convex combinations of Nash equilibria
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A Correlated Equilibrium is a Nash equilibrium when the players receive payoff-irrelevant signals before playing the game (Aumann 1974)

- Independent signals Nash equilibrium
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 convex combinations of Nash equilibria
- Butterflies play the Chicken Game ("Speckled Wood" Pararge aegeria)
- Boston Celtics' front line

Correlated Equilibrium

Signals (public, correlated) are unavoidable

Correlated Equilibrium

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- Common Knowledge of Rationality ⇔ Correlated Equilibrium (Aumann 1987)

Correlated Equilibrium

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A joint distribution z is a correlated equilibrium

$$\Leftrightarrow$$

$$\sum_{s^{-i}} u(j, s^{-i}) z(j, s^{-i}) \geq \sum_{s^{-i}} u(k, s^{-i}) z(j, s^{-i})$$

for all $i \in N$ and all $j,k \in S^i$

RESULT

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THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

Regret Matching

Hart and Mas-Colell, Ec'ca 2000

RESULT

THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

- Regret Matching
- General regret-based dynamics

Hart and Mas-Colell, Ec'ca 2000, JET 2001

Regret Matching

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Regret Matching

"REGRET": the increase in past payoff, if any, if a different action would have been used

"MATCHING": switching to a different action with a probability that is proportional to the regret for that action

THERE EXIST general, natural dynamics leading to CORRELATED EQUILIBRIA

"general": in all games

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 - adaptive (also: close to "behavioral")

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- "general": in all games
- "natural":
 - adaptive (also: close to "behavioral")
 - simple and efficient: computation, time, information
- "leading to correlated equilibria": statistics of play become close to CORRELATED EQUILIBRIA

NASH EQUILIBRIUM: a fixed-point of a non-linear map

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- CORRELATED EQUILIBRIUM: a solution of finitely many linear inequalities

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set-valued fixed-point (curb sets)?

"LAW OF CONSERVATION OF COORDINATION":

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There must be some coordination —

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There must be some COORDINATION—
either in the EQUILIBRIUM notion,

"LAW OF CONSERVATION OF COORDINATION":

There must be some COORDINATION—
either in the EQUILIBRIUM notion,
or in the DYNAMIC

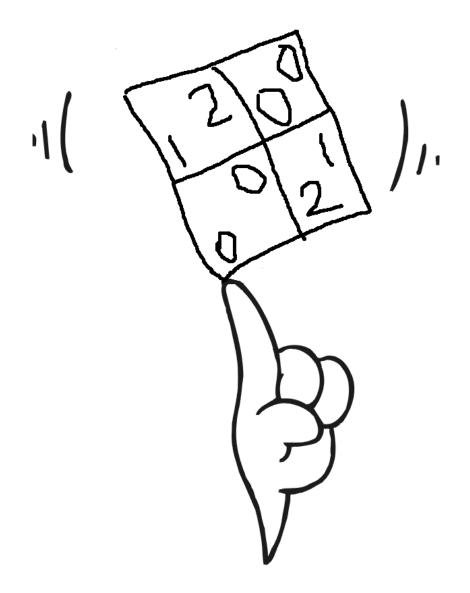
A. Demarcate the **BORDER** between

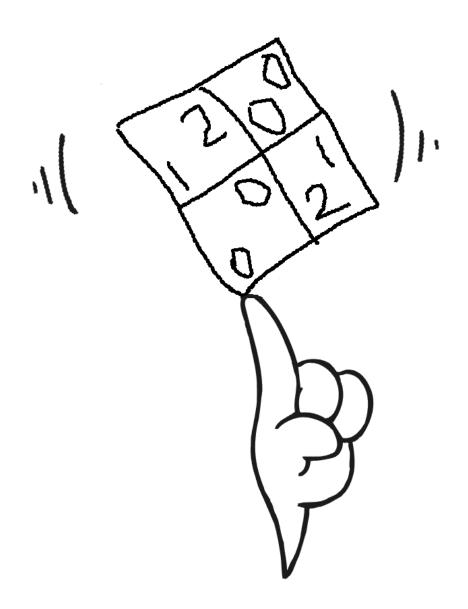
- A. Demarcate the **BORDER** between
 - classes of dynamics where convergence to equilibria CAN be obtained

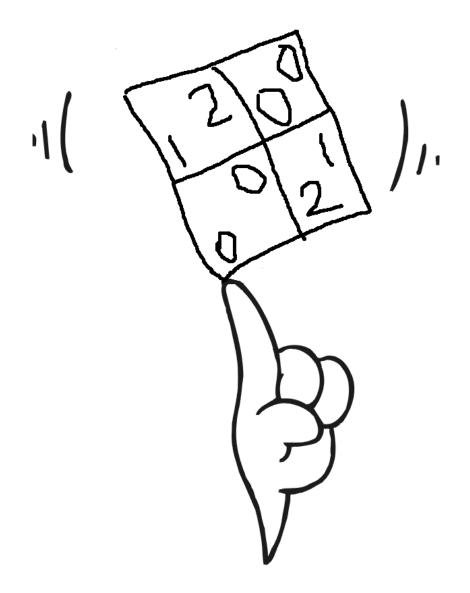
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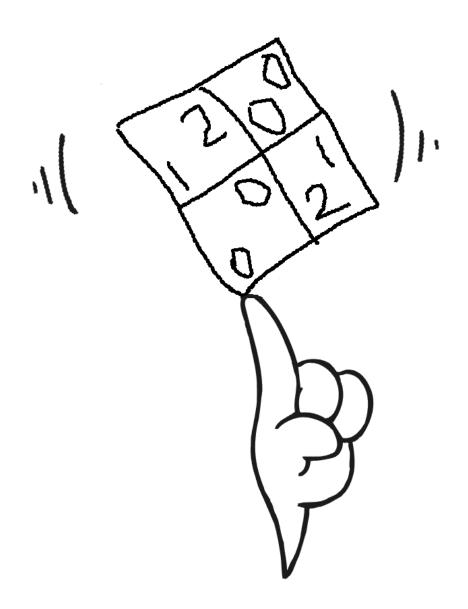
- classes of dynamics where convergence to equilibria CAN be obtained, and
- classes of dynamics where convergence to equilibria CANNOT be obtained

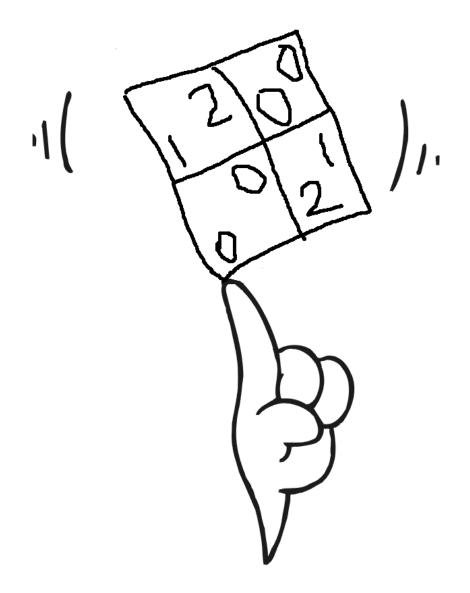
- A. Demarcate the **BORDER** between
 - classes of dynamics where convergence to equilibria
 CAN be obtained, and
 - classes of dynamics where convergence to equilibria CANNOT be obtained
- B. Find NATURAL dynamics for the various equilibrium concepts

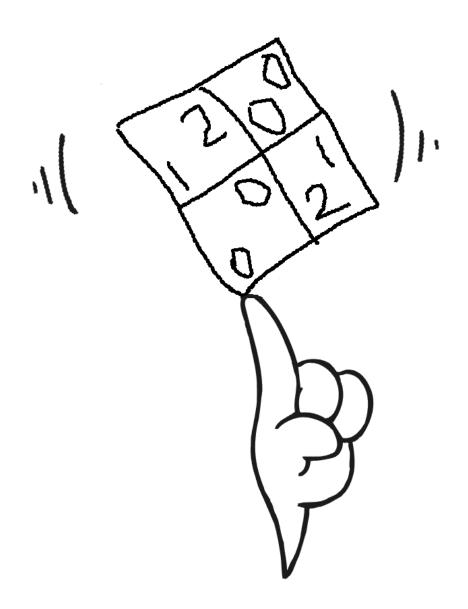


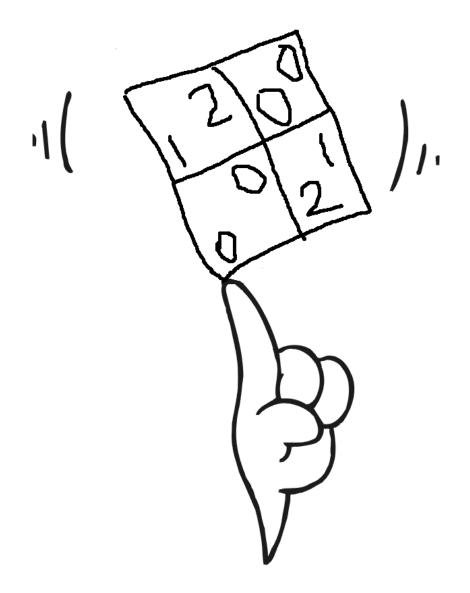


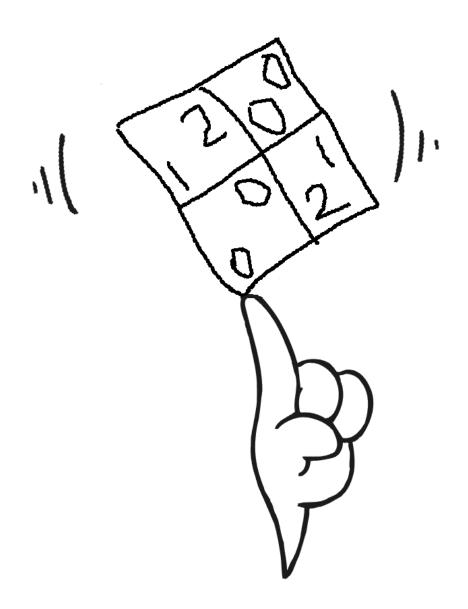




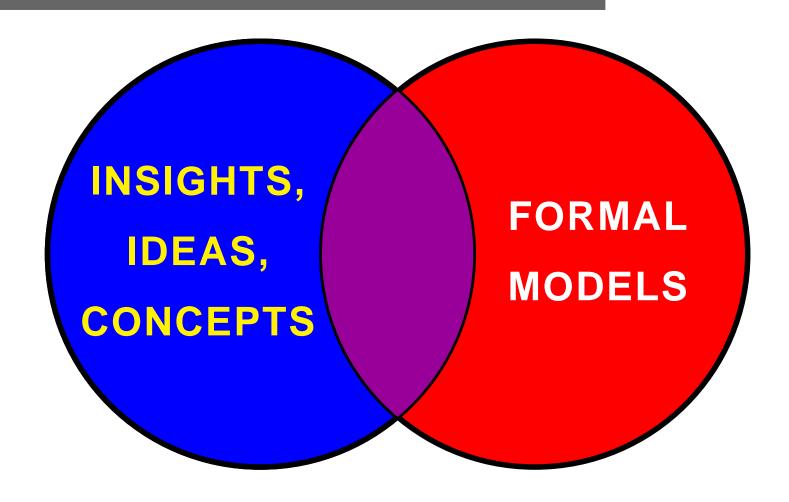












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